

## Linear and Integer Optimization

### Exercise Sheet 9

**Exercise 9.1:** Show that in the inequality

$$\max\{c^t x \mid Ax \leq b, x \in \mathbb{Z}^n\} \leq \min\{b^t y \mid A^t y = c, y \geq 0, y \in \mathbb{Z}^m\}$$

in general, equality does not hold, even if the two corresponding optimization problems are feasible and bounded. (2 Points)

**Exercise 9.2:** Give an example each of

1. a full-dimensional unbounded rational polyhedron  $P$  such that  $P_I$  is empty.
2. an unbounded polyhedron  $P$  such that  $P_I$  is non-empty and bounded.
3. a polyhedron  $P$  such that  $P_I \neq \emptyset$  is not closed.
4. a feasible and bounded ILP without optimum solution. (2+2+2+2 Points)

**Exercise 9.3:** Let  $P, Q \subseteq \mathbb{R}^n$  be two polyhedra. Show that  $P_I + Q_I \subseteq (P + Q)_I$ . Give an example where  $P_I + Q_I \neq (P + Q)_I$ . (5 points)

**Exercise 9.4:** The MAXIMUM-STABLE-SET-PROBLEM is defined as follows. Given a graph  $G$ , we are looking for a vertex set  $S \subseteq V(G)$  of maximum cardinality  $|S|$  such that  $\{v, w\} \notin E(G)$  for all  $v, w \in S$ . A vertex set  $S$  with  $E(G[S]) = \emptyset$  is called *stable set*. This problem can be modeled by the following ILP:

$$\max \sum_{v \in V(G)} x_v \tag{1}$$

$$s.d. \quad x_v + x_w \leq 1 \quad \forall \{v, w\} \in E(G) \tag{2}$$

$$x_v \in \{0, 1\} \quad \forall v \in V(G) \tag{3}$$

1. Prove that the integrality gap of the LP relaxation is at least  $\frac{|V(G)|}{2}$ .
2. Show that following inequalities are valid for the MAXIMUM-STABLE-SET-PROBLEM

$$\sum_{v \in V(H)} x_v \leq \left\lceil \frac{|V(H)| - 1}{2} \right\rceil \quad \text{for a circuit } H \subseteq G.$$

Give examples, where the optima of the linear relaxations are reduced (strictly). (2+3 Points)

**Submission deadline:** Tuesday, June 23, 2026, 16:00, via eCampus (in groups of at most 3 students).