

## Programming Exercise 2

**Exercise P.2.** Let  $G = (V, E)$  be an undirected graph and  $k = |V|$ . Let  $C \subseteq V$  and  $f : V \setminus C \rightarrow \{1, \dots, k\}$  be a placement function (in particular  $f$  is injective).

- (a) Compute positions  $g : V \rightarrow [1, k]$ , extending  $f$ , approx. minimizing

$$\sum_{e=\{v,w\} \in E} |g(v) - g(w)|^2$$

This can be done by choosing any initial positions and then iteratively picking the average position of all neighbors for each circuit (this is equivalent to perform Gauß-Seidel iterations on the linear equation system defined in Proposition 4.18 in the lecture). Stop if the maximum circuit movement in one iteration is below 0.1. Note that  $g$  is not required to be injective.

- (b) Use Nesterov's Accelerated Gradient Method to find positions  $g : V \rightarrow [1, k]$ , extending  $f$ , to approx. minimize the weighted average netlength

$$\text{NWA}_\gamma(f) := \sum_{\{v,w\} \in E} \text{WA}_\gamma(f(v), f(w)) - \text{WA}_\gamma(-f(v), -f(w)).$$

Again,  $g$  is not required to be injective.

- (c) Finally extend the placement  $f$  to a placement  $f : V \rightarrow \{1, \dots, k\}$ . First use the program from the first, then from the second task obtaining some  $g$  which is not necessarily injective and integral. Sort the circuits by their position w.r.t.  $g$ , pick the median unassigned circuit and assign it to the median available position in  $f$ . Assign all the other circuits to either the left or the right side depending on their positions w.r.t.  $g$  and recursively solve these two smaller instances, finally yielding a placement  $f$  that is integral and distinct.

Run your programs on the instances on the website. For each of the tasks output the linear length and the sum of the quadratic lengths. Moreover print the positions  $g$  of the circuits  $C$  for the first two tasks and the positions

$f$  for all vertices for the third task, by printing a single line for each circuit containing its index in the input and its computed position.

The instances are given in DIMACS format:

- The first line starts with a p (problem) and has the following format:  
p edge  $k$   $m$   
where  $k$  = number of vertices  $(1, \dots, k)$  and  $m$  = number of edges  $(1, \dots, m)$ .
- Lines starting with an e define edges and have the following format:  
e  $i$   $j$   
where the edge is joining the vertices indexed by  $i$  and  $j$ .
- Lines starting with an n define vertex-positions and have the following format:  
n  $i$   $p$   
where  $i$  is the vertex index and the integer  $p \in \{1, \dots, k\}$  is its position. Vertices without fixed position, for which you should compute a position, will have  $p = -1$ .

There is a C-routine for reading in a graph in the given format provided on the website.

The program must be written in C or C++ and must compile and run on Linux. You are allowed to use any any ISO C or C++ standard including C++20. You can use any tool available in the standard library. Your program must compile with either Clang (any version  $\geq 3.4.2$ ) or Gcc (version  $\geq 4.8.3$ ) with `-Wall -Wextra -Wpedantic -Werror` and cannot link to any external library except the standard library. To achieve the maximum score, your program must not leak any memory and must be well documented.

(12+12+12 points)

**Deadline:** June 25<sup>th</sup>, via email to schlomberg@or.uni-bonn.de. The websites for the lecture with all exercises and test instances can be found at:

[http://www.or.uni-bonn.de/lectures/ss24/chipss24\\_ex.html](http://www.or.uni-bonn.de/lectures/ss24/chipss24_ex.html)