

## Linear and Integer Optimization

### Assignment Sheet 9

1. Let  $G$  be simple undirected graph. Consider the following linear program:

$$\begin{aligned}
 \min \quad & \sum_{e=\{v,w\} \in E(G)} x_{vw} \\
 \text{s.t.} \quad & \sum_{w \in S} x_{vw} \geq \lceil \frac{1}{4}|S|^2 + \frac{1}{2}|S| \rceil \quad \text{for } v \in V(G), S \subseteq V(G) \setminus \{v\} \\
 & x_{uw} \leq x_{uv} + x_{vw} \quad \text{for } u, v, w \in V(G) \\
 & x_{vw} \geq 0 \quad \text{for } v \in V(G) \\
 & x_{vv} = 0 \quad \text{for } v \in V(G)
 \end{aligned}$$

- (a) Show that this is an LP-relaxation of the following problem: Find distances  $x_{vw}$  for the nodes of  $G$  such that  $\sum_{e=\{v,w\} \in E(G)} x_{vw}$  is minimized under the condition that there is an ordering  $\{v_1, \dots, v_{|V(G)|}\} = V(G)$  with  $x_{v_i v_j} = |i - j|$  for  $i, j \in \{1, \dots, |V(G)|\}$ .
- (b) Prove that there is a polynomial-time separation oracle for the polyhedron of the feasible solutions of the LP. (4+4 points)

2. Consider the following optimization problem:

$$\begin{aligned}
 \min \quad & \frac{c^t x + d}{f^t x + g} \\
 \text{s.t.} \quad & Ax \leq b \\
 & \|x\|_\infty \leq R
 \end{aligned}$$

where  $c, f \in \mathbb{Q}^n$ ,  $d, g, R \in \mathbb{Q}$ ,  $A \in \mathbb{Q}^{m \times n}$ ,  $b \in \mathbb{Q}^m$ . You may assume that  $f^t x + g > 0$  and  $c^t x + d > 0$  for any  $x \in \mathbb{R}^n$  with  $\|x\|_\infty \leq R$  and that there is a feasible solution. Show that for any  $\epsilon > 0$  there is a polynomial-time algorithm computing a feasible solution  $x^*$  with  $\frac{c^t x^* + d}{f^t x^* + g} \leq \text{OPT}(1 + \epsilon)$  where OPT is the value of an optimum solution. (6 points)

3. Consider the following situation: You are a foreign exchange trader trading in  $n$  different currencies with an initial budget  $b = (b_1, \dots, b_n) \geq 0$ , where  $b_i$  is the amount of currency  $i$ . For every pair  $1 \leq i, j \leq n$  of currencies, you are given an exchange rate  $r_{ij}$ . Assume that you are able to exchange an arbitrary (non-negative) amount  $x_{ij}$  of currency  $i$  into  $r_{ij}x_{ij}$  units of currency  $j$ . You want to determine trades that will maximize the amount of currency 1 available to you in the end. You may assume that all trades can occur simultaneously (e.g. by cost-free borrowing for the instant during which the trades occur), so long as the final amounts in all currencies are non-negative.

Model this problem as a linear program and derive its dual. (5 points)

*Remark:* If the LP is unbounded, one speaks of an *arbitrage opportunity*, i.e. a situation that allows “risk-free” profit.

4. A parcel service leases cars on a basis of 3, 4 or 5 months. The cost for a 3-months leasing contract (for one car) is 1700 EUR, for a 4-months contract 2200 EUR, and for 5-months contract 2500 EUR. For a certain period (e.g. one year), the company knows in advance for each month how many vehicles will be needed in that month. Formulate the problem of finding a cheapest way to lease sufficiently many cars as a linear program. Show in particular that the linear program always has an optimum solution that is integral. (6 points)