

## Combinatorial Optimization

### Exercise Sheet 8

**Exercise 8.1:** Given an undirected graph, an *odd cycle cover* is defined to be a subset of the edges containing at least one edge of each odd circuit. Show how to find in polynomial time a minimum weight odd cycle cover in a planar graph with nonnegative weights on the edges. Can you also solve the problem for general weights?

(4 Points)

**Exercise 8.2:** Let  $G$  be an undirected graph and  $T \subseteq V(G)$  with  $|T| = 2k$  even. Prove that the minimum cardinality of a  $T$ -cut in  $G$  equals the maximum of  $\min_{i=1}^k \lambda_{s_i, t_i}$  over all pairings  $T = \{s_1, t_1, \dots, s_k, t_k\}$ , where  $\lambda_{s,t}$  denotes the maximum number of pairwise edge-disjoint  $s$ - $t$ -paths.

Hint: Use the Padberg-Rao-Theorem.

(4 Points)

**Exercise 8.3:** The DIRECTED CHINESE POSTMAN PROBLEM can be formulated as follows: Given a strongly connected simple digraph  $G$  with weights  $c : E(G) \rightarrow \mathbb{R}_+$ , find  $f : E(G) \rightarrow \mathbb{N}$  such that the graph which contains  $f(e)$  copies of each edge  $e \in E(G)$  is Eulerian and  $\sum_{e \in E(G)} c(e)f(e)$  is minimum. Show how to solve this problem in polynomial time by reducing it to a MINIMUM COST FLOW PROBLEM.

(4 Points)

**Exercise 8.4:** Let  $G$  be an undirected graph and  $b_1, b_2 : V(G) \rightarrow \mathbb{N}$ . Describe the convex hull of functions  $f : E(G) \rightarrow \mathbb{Z}_+$  with  $b_1(v) \leq \sum_{e \in \delta(v)} f(e) \leq b_2(v)$ .

Hint: For  $X, Y \subseteq V(G)$  with  $X \cap Y = \emptyset$  consider the constraint

$$\sum_{e \in E(G[X])} f(e) - \sum_{e \in E(G[Y]) \cup E(Y, Z)} f(e) \leq \left\lfloor \frac{1}{2} \left( \sum_{x \in X} b_2(x) - \sum_{y \in Y} b_1(y) \right) \right\rfloor,$$

where  $Z := V(G) \setminus (X \cup Y)$ . Use the b-matching polytope.

(4 Points)

**Deadline:** Tuesday, December 22, 2015, **before** the lecture.

**Information:** Submissions by groups of up to **three** students are allowed.