## Combinatorial Optimization

## Exercise Sheet 6

Exercise 6.1: Prove:

- (i) For every two-edge-connected graph G and every  $T \subseteq V(G)$  with |T| even there exists a T-join J in G with  $|J| \leq \frac{|E(G)|}{2}$ . (2 Points)
- (ii) For every factor-critical graph G and every  $T \subseteq V(G)$  with |T| even there exists a T-join J in G with  $|J| \leq \frac{|V(G)|-1}{2}$ . (2 Points)

**Exercise 6.2:** For  $n \in \mathbb{N}$ , let  $P_n$  be the convex hull of all even 0-1-vectors. More precisely, let

$$P_n = \operatorname{conv} \{ x \in \{0, 1\}^n : \sum_{i=0}^n x_i \equiv 0 \pmod{2} \}.$$

Prove that the relaxation complexity is of size  $rc(P_n) = 2^{\Theta(n)}$ . (4 Points)

**Exercise 6.3:** Let G = (V, E) be an undirected graph and n := |V|. Prove that the following inequality system with  $\mathcal{O}(n^3)$  variables and constraints describes a polytope whose orthogonal projection to the *x*-variables yields the spanning tree polytope of G.

$$x_{e} \geq 0 \qquad (e \in E)$$

$$z_{u,v,w} \geq 0 \qquad (\{u,v\} \in E, w \in V \setminus \{u,v\})$$

$$\sum_{e \in E} x_{e} = n - 1$$

$$x_{e} = z_{u,v,w} + z_{v,u,w} \qquad (e = \{u,v\} \in E, w \in V \setminus e)$$

$$x_{e} + \sum_{\{u,v\} \in \delta(v) \setminus \{e\}} z_{u,v,w} = 1 \qquad (v \in V, e = \{v,w\} \in \delta(v))$$

(4 Points)

**Exercise 6.4:** Show that the following algorithm finds in a graph G (which is not a forest) with edge weights  $w : E(G) \to \mathbb{R}$  a cycle  $C \subset E(G)$  that minimizes  $\frac{w(C)}{|C|}$  in strongly polynomial time: First reduce all edge weights by  $\max\{w(e)|e \in E(G)\}$ . Then find a minimum-weight  $\emptyset$ -join J. If w(J) = 0 output a cycle of weight 0, otherwise add  $\frac{-w(J)}{|J|}$  to all edge weights and iterate (i.e. find again a minimum-weight  $\emptyset$ -join). (4 Points)

**Deadline:** Tuesday, December 8, 2015, **before** the lecture. **Information:** Submissions by groups of up to **three** students are allowed.