

Combinatorial Optimization

Exercise Sheet 3

Exercise 3.1:

Prove:

- (i) A minimal factor-critical graph G has at most $\frac{3}{2}(|V(G)| - 1)$ edges and that this bound is tight. (2 Points)
- (ii) Let G be a graph and M a matching in G . If $X \subseteq V(G)$ is the set of M -exposed vertices, then a shortest M -alternating X - X -walk of positive length can be found in $O(|E(G)|)$ time. (2 Points)

Exercise 3.2:

Prove that an undirected graph G is factor-critical if and only if G is connected and $\nu(G) = \nu(G - v)$ for all $v \in V(G)$. (3 Points)

Exercise 3.3:

Let G be a graph and M a matching in G that is not maximum.

- (i) Show that there are $\nu(G) - |M|$ vertex-disjoint M -augmenting paths in G .
- (ii) Prove that there exists an M -augmenting path of length at most $\frac{\nu(G)+|M|}{\nu(G)-|M|}$.
- (iii) Let P be a shortest M -augmenting path in G and P' an $(M \triangle E(P))$ -augmenting path. Prove $|E(P')| \geq |E(P)| + 2|E(P \cap P')|$.

Consider the following algorithm: We start with the empty matching and in each iteration augment the matching along a shortest augmenting path. Let P_1, P_2, \dots be the sequence of augmenting paths chosen.

- (iv) Show that if $|E(P_i)| = |E(P_j)|$ for $i \neq j$, then P_i and P_j are vertex-disjoint.
- (v) Conclude that the sequence $|E(P_1)|, |E(P_2)|, \dots$ contains at most $2\sqrt{\nu(G)} + 2$ different numbers.

(5 Points)

Exercise 3.4:

Let $G = (V, E)$ a graph and $X \subseteq V$. Let $\beta(G, X)$ be the maximum size of a set $Y \subseteq X$ for which there is a matching in G that covers Y . Prove

$$\beta(G, X) = \min_{U \subseteq V} |X| + |U| - q_X(U).$$

Here $q_X(U)$ denotes the number of odd connected components of $G-U$ whose vertices are all in X .

Hint: Construct a new graph with $2|V|$ vertices and apply Tutte's Theorem. (4 Points)

Deadline: Tuesday, November 17, 2015, **before** the lecture.

Information: Submissions by groups of up to **three** students are allowed.