Exercise 3.1:
Prove:

(i) A minimal factor-critical graph $G$ has at most $\frac{3}{2}(|V(G)| - 1)$ edges and that this bound is tight. (2 Points)

(ii) Let $G$ be a graph and $M$ a matching in $G$. If $X \subseteq V(G)$ is the set of $M$-exposed vertices, then a shortest $M$-alternating $X$-$X$-walk of positive length can be found in $O(|E(G)|)$ time. (2 Points)

Exercise 3.2:
Prove that an undirected graph $G$ is factor-critical if and only if $G$ is connected and $\nu(G) = \nu(G - v)$ for all $v \in V(G)$. (3 Points)

Exercise 3.3:
Let $G$ be a graph and $M$ a matching in $G$ that is not maximum.

(i) Show that there are $\nu(G) - |M|$ vertex-disjoint $M$-augmenting paths in $G$.

(ii) Prove that there exists an $M$-augmenting path of length at most $\frac{\nu(G) + |M|}{\nu(G) - |M|}$.

(iii) Let $P$ be a shortest $M$-augmenting path in $G$ and $P'$ an $(M \triangle E(P))$-augmenting path. Prove $|E(P')| \geq |E(P)| + 2|E(P \cap P')|$.

Consider the following algorithm: We start with the empty matching and in each iteration augment the matching along a shortest augmenting path. Let $P_1, P_2, \ldots$ be the sequence of augmenting paths chosen.

(iv) Show that if $|E(P_i)| = |E(P_j)|$ for $i \neq j$, then $P_i$ and $P_j$ are vertex-disjoint.

(v) Conclude that the sequence $|E(P_1)|, |E(P_2)|, \ldots$ contains at most $2\sqrt{\nu(G)} + 2$ different numbers. (5 Points)
Exercise 3.4:
Let $G = (V, E)$ a graph and $X \subseteq V$. Let $\beta(G, X)$ be the maximum size of a set $Y \subseteq X$ for which there is a matching in $G$ that covers $Y$. Prove

$$\beta(G, X) = \min_{U \subseteq V} |X| + |U| - q_X(U).$$

Here $q_X(U)$ denotes the number of odd connected components of $G - U$ whose vertices are all in $X$.

*Hint:* Construct a new graph with $2|V|$ vertices and apply Tutte’s Theorem. (4 Points)

**Deadline:** Tuesday, November 17, 2015, before the lecture.

**Information:** Submissions by groups of up to three students are allowed.