## Exercise Set 9

**Exercise 9.1.** Consider the PLACEMENT LEGALIZATION PROBLEM (see lecture notes for the formal definition) with  $y_{\text{max}} - y_{\text{min}} = 1$ . We are given an infeasible placement  $\tilde{x} : \mathcal{C} \to \mathbb{R}$ . Show that there are feasible instances for which there is no optimum solution which is consistent with  $\tilde{x}$ , i.e. such that  $x(C) < x(C') \Rightarrow \tilde{x}(C) \leq \tilde{x}(C')$ .

(3 points)

**Exercise 9.2.** Consider following specialized legalization problem. A feasible placement consists of integral placement coordinates  $x, y : \mathcal{C} \to \mathbb{Z}^2$  (such that all circuits are located within the chip image), and all circuits have unit height and width. In addition, there is an (infeasible) input placement  $\tilde{x}, \tilde{y} : \mathcal{C} \to \mathbb{R}^2$ .

- Find a polynomial-time algorithm that finds a legal placement minimizing the linear or quadratic movement.
- Find a linear time algorithm that finds a legal placement minimizing the quadratic movement if the instance consists of a single row. Assume that the input is sorted.

(2+2 points)

**Exercise 9.3.** Consider the following variant of the SINGLE ROW PLACE-MENT WITH FIXED ORDERING problem, in which we minimize the bounding box net length:

**Input:** A set  $C = \{C_1, \ldots, C_n\}$  of circuits, widths  $w(C_i) \in \mathbb{R}_+$ , an interval  $[0, w(\Box)]$ , s.t.  $\sum_{i=1}^n w(C_i) \leq w(\Box)$ . A netlist  $(C, P, \gamma, \mathcal{N})$ where the offset of a pin  $p \in P$  satisfies  $x(p) \in [0, w(\gamma(p))]$ . Weights  $\alpha : \mathcal{N} \to \mathbb{R}_+$ .

**Task:** Find a feasible placement given by a function  $x : \mathcal{C} \to \mathbb{R}$ s.t.  $0 \leq x(C_1), x(C_i) + w(C_i) \leq x(C_{i+1})$  for  $i = 1, \ldots, n-1$  and  $x(C_n) + w(C_n) \leq w(\Box)$ , that minimizes

$$\sum_{N \in \mathcal{N}} \alpha(N) \cdot \mathrm{BB}(N).$$

Here, BB(N) denotes the bounding box net length.

Show that there exist  $f_i : [0, w(\Box)] \to \mathbb{R}, i = 1, ..., n$ , piecewise linear, continuous and convex, such that we can solve this problem by means of the SINGLE ROW ALGORITHM.

(5 points)

**Deadline:** June 13, before the lecture. The websites for lecture and exercises can be found at:

## http://www.or.uni-bonn.de/lectures/ss23/chipss23\_ex.html

In case of any questions feel free to contact me at drees@or.uni-bonn.de.