

## Exercise Set 3

**Exercise 3.1.** Let  $T$  be an instance of the Rectilinear Steiner Tree Problem and  $r \in T$ . For a rectilinear Steiner tree  $Y$  we denote by  $f(Y)$  the maximum length of a path from  $r$  to any element of  $T \setminus \{r\}$  in  $Y$ .

- (a) Find an instance where no Steiner tree minimizes both length and  $f$ .
- (b) Consider the problem of finding a shortest Steiner tree  $Y$  minimizing  $f(Y)$  among all shortest Steiner trees. Is there always a tree with these properties which is a subgraph of the Hanan grid?

(1 + 3 points)

**Exercise 3.2.** Let  $Y$  be a Steiner tree for terminal set  $T$  with  $|T| \geq 2$  in which all leaves are terminals. Prove

$$\sum_{t \in T} (|\delta_Y(t)| - 1) = k - 1$$

where  $k$  is the number of full components of  $Y$ .

(3 points)

**Exercise 3.3.** Let  $T$  be a finite set of terminals.

A topology for  $T$  is a tree  $X$  with  $T \subseteq V(X)$ , such that all Steiner points (elements of  $V(X) \setminus T$ ) have degree at least 3. A topology is a standard topology if all Steiner points have degree exactly 3 in  $X$  and all elements of  $T$  are leaves.

We say that a Steiner tree  $Y$  has topology  $X$  if there is a map  $\varphi: V(X) \rightarrow V(Y)$  with  $\varphi(t) = t$  for all  $t \in T$  such that  $\{E(P_e) : e \in E(X)\}$  is a partition of  $E(Y)$ , where for  $e = \{v, w\} \in E(X)$ ,  $P_e$  is the path from  $\varphi(v)$  to  $\varphi(w)$  in  $Y$ .

Modify the DIJKSTRA-STEINER Algorithm such that it finds an optimum Steiner tree with a fixed standard topology in  $\mathcal{O}(k(m + n \log n))$  time, where  $n = |V(G)|$ ,  $m = |E(G)|$  and  $k = |T|$ .

(5 points)

**Deadline:** April 25, before the lecture. The websites for lecture and exercises can be found at:

`http://www.or.uni-bonn.de/lectures/ss23/chipss23\_ex.html`

In case of any questions feel free to contact me at `drees@or.uni-bonn.de`.