

## Exercise Set 11

**Exercise 11.1.** Let  $G = (V, E)$  be an undirected graph with non-negative weights  $w : V \rightarrow \mathbb{R}$ , a set of sinks  $T \subset V$ , and a root vertex  $r \in V \setminus T$ . Additionally, we are given required arrival times  $rat : T \rightarrow \mathbb{R}$ . The goal of the DELAY BOUNDED STEINER TREE PROBLEM is to compute a steiner tree  $S$  of  $\{r\} \cup T$  in  $G$  with minimum weight, such that for each  $t \in T$  the length of the unique  $r$ - $t$  path in  $S$  is at most  $rat(t)$ . Assuming  $P \neq NP$ , show that there is no  $\log(|T|)$  approximation algorithm for this problem.

*Hint:* You may use that it is NP hard to find a  $\log(n)$  approximation for SET COVER.

(5 points)

**Exercise 11.2.** Let  $t_1, \dots, t_n \in \mathbb{R}_{>0}^2$ ,  $r := (0, 0) \in \mathbb{R}^2$ ,  $d(x, y) := \|x - y\|_1$ .

- (a) Show that there exists a perfect matching on  $t_1, \dots, t_n$  with length at most that of a steiner arborescence on  $t_1, \dots, t_n$  rooted in  $r$ .
- (b) Describe a polynomial time algorithm that computes a  $\mathcal{O}(\log(n))$  approximation for a minimum length Steiner arborescence on  $t_1, \dots, t_n$  rooted in  $r$ , such that the length of each  $r$ - $t_i$  path is  $\|r - t_i\|_1$  ( $i = 1, \dots, n$ ).

(2+3 points)

**Exercise 11.3.** Let  $\alpha > 1$  and  $1 \leq \beta < 1 + 2/(\alpha - 1)$ . Construct a connected, planar graph  $G$  with  $w : E(G) \rightarrow \mathbb{R}_+$  and  $r \in V(G)$  that contains no spanning tree  $T$  with the following properties:

- (a) For each  $v \in V(G)$ :  $\text{dist}_{w,T}(r, v) \leq \alpha \cdot \text{dist}_{w,G}(r, v)$ .
- (b) For a minimum-spanning tree  $M$ :  $\sum_{e \in E(T)} w(e) \leq \beta \cdot \sum_{e \in E(M)} w(e)$ .

(5 points)

**Exercise 11.4.** Given a power consumption  $P_l > 0$  for each buffer  $l \in L$ , extend the algorithm by van Ginneken from the lecture to obtain a PTAS for the problem of finding an assignment of buffers such that all required arrival times are met and power is minimized.

- (a) First assume that for all  $l \in L$ , both  $P_l$  and  $\frac{1}{P_l}$  are polynomially bounded in the input size.
- (b) Then solve the general case using binary search.

(3+2 points)

**Deadline:** Tuesday, July 2<sup>nd</sup>, before the lecture. The websites for lecture and exercises can be found at:

<http://www.or.uni-bonn.de/lectures/ss19/chipss19.html>

In case of any questions feel free to contact me at [klotz@or.uni-bonn.de](mailto:klotz@or.uni-bonn.de).