Exercise Set 6

Exercise 6.1. Consider the spreading LP for d = 2:

$$\begin{array}{ll} \min & \sum_{e \in E(G)} w(e) \, l(e) \\ \text{s.t.} & \sum_{y \in X} l(\{x, y\}) \geq \frac{1}{4} \left(|X| - 1 \right)^{3/2} & x \in X \subseteq V(G) \\ & l(\{x, y\}) + l(\{y, z\}) \geq l(\{x, z\}) & x, y, z \in V(G) \\ & l(\{x, y\}) \geq 1 & x, y \in V(G), \; x \neq y \\ & l(\{x, x\}) = 0 & x \in V(G) \end{array}$$

Show that the optimum of the spreading LP is a lower bound for the cost of any 2-dimensional arrangement.

(5 points)

Exercise 6.2. Provide a polynomial time algorithm for the STANDARD PLACEMENT PROBLEM restricted to instances with only one circuit.

(5 points)

Exercise 6.3. Consider the STANDARD PLACEMENT PROBLEM on instances without blockages, where $h(C) \equiv 1 \equiv w(C)$ (unit size for $C \in C$) as well as $w(N) \equiv 1$ (unit net weights for $N \in \mathcal{N}$).

Prove or disprove that this problem is NP-hard.

(5 points)

Exercise 6.4. Let G = (V, E) be an undirected graph with edge weights $w : E \to \mathbb{R}_{\geq 0}$ and $k \in \mathbb{N}$. Let $C \subseteq V$ and $f : V \setminus C \to \{1, \ldots, k\}$ be a placement function. We are looking for positions $f : C \to \{1, \ldots, k\}$ s.t.

$$\sum_{e=\{v,w\}\in E} w(e)\cdot |f(v) - f(w)|$$

is minimum. Note that f is not required to be injective.

Prove that this problem can be solved optimally by solving k-1 minimum weight *s*-*t*-cut problems in digraphs with $\mathcal{O}(|V|)$ vertices and $\mathcal{O}(|E|)$ edges.

Hint: Consider digraphs $G_j = (V_j, E_j)$ with $V_j := \{s, t\} \cup C$ and

$$E_j := \left\{ (s, v) : \exists w \in V \setminus C, f(w) \leq j, \{v, w\} \in E \right\} \cup \\ \left\{ (v, w) : v, w \in C, \{v, w\} \in E \right\} \cup \\ \left\{ (v, t) : \exists w \in V \setminus C, f(w) > j, \{v, w\} \in E \right\}$$

(5 points)

Deadline: May 16th, before the lecture. The websites for lecture and exercises can be found at:

In case of any questions feel free to contact me at klotz@or.uni-bonn.de.