

## Exercise Set 5

**Exercise 5.1.** Consider the following CLUSTERED RECTILINEAR STEINER TREE PROBLEM: Given a partition  $T = \dot{\bigcup}_{i=1}^k P_i$  of the terminals ( $\emptyset \neq P_i \subseteq \mathbb{R}^2, |P_i| < \infty$ ), find a (rectilinear) Steiner tree  $Y_i$  for each set of terminals  $P_i$  and one rectilinear, toplevel (group) Steiner tree  $Y_{\text{top}}$  connecting the embedded trees  $Y_i$  ( $i = 1, \dots, k$ ). The task is to minimize the total length of all trees.

Let  $A$  be an  $\alpha$ -approximation algorithm for the RECTILINEAR STEINER TREE PROBLEM. A feasible solution to the CLUSTERED RECTILINEAR STEINER TREE PROBLEM can be found by first selecting a connection point  $q_i \in \mathbb{R}^2$  for each  $i = 1, \dots, k$  and then computing  $Y_i := A(P_i \cup \{q_i\})$  and  $Y_{\text{top}} := A(\{q_i : 1 \leq i \leq k\})$ .

- (a) Show that picking  $q_i \in P_i$  arbitrarily yields a  $2\alpha$  approximation.
- (b) Prove that choosing each  $q_i$  as the center of the bounding box of  $P_i$  implies a  $\frac{7}{4}\alpha$  approximation algorithm.
- (c) Show that both approximation ratios above are tight.

(2 + 4 + 2 points)

**Exercise 5.2.** Prove that the STANDARD PLACEMENT PROBLEM can be solved optimally in

$$O\left(\left((n+s)!\right)^2 \left((m+n^2+k \log k)(n+k) \log(n+k) + (sn)\right)\right)$$

time, where  $n := |\mathcal{C}|$ ,  $k := |\mathcal{N}|$ ,  $m := |\mathcal{P}|$  and  $s := |\mathcal{S}|$ .

(6 points)

**Exercise 5.3.** Consider quadratic netlength minimization in  $x$ -dimension based on the (quadratic) CLIQUE netmodel i.e.

$$\text{CLIQESQ}(N) := \sum_{\{p,q\} \subseteq N} \frac{w(N)}{|N| - 1} \left( x(p) + x(\gamma(p)) - x(q) - x(\gamma(q)) \right)^2$$

- (a) Show that CLIQUESQ can be replaced equivalently by the quadratic STARSQ netmodel

$$\text{STARSQ}(N) := w'(N) \cdot \min \left\{ \sum_{p \in N} (x(p) + x(\gamma(p)) - c)^2 \mid c \in \mathbb{R} \right\}$$

for an appropriate weight function  $w'$ .

- (b) For a fixed placement  $x$  and a single net  $N$  let  $l, r \in N$  be defined as  $l := \arg \min \{x(p) + x(\gamma(p)) \mid p \in N\}$  and  $r := \arg \max \{x(p) + x(\gamma(p)) \mid p \in N\}$ . We further define for  $p, q \in N$

$$w_{pq}^{\text{B2B}} := \begin{cases} 0 & \text{if } \{p, q\} \cap \{l, r\} = \emptyset, \\ \left| x(q) + x(\gamma(q)) - x(p) - x(\gamma(p)) \right|^{-1} & \text{else.} \end{cases}$$

Show that the CLIQUESQ netlength with weights  $w^{\text{B2B}}$  equals the (linear) bounding box netlength for placement  $x$ .

(3 + 3 points)

**Deadline:** May 9<sup>th</sup>, before the lecture. The websites for lecture and exercises can be found at:

<http://www.or.uni-bonn.de/lectures/ss19/chipss19.html>

In case of any questions feel free to contact me at [klotz@or.uni-bonn.de](mailto:klotz@or.uni-bonn.de).