

## Exercise Set 4

**Exercise 4.1.** For a finite set  $\emptyset \neq T \subsetneq \mathbb{R}^2$  we define

$$\text{BB}(T) := \max_{(x,y) \in T} x - \min_{(x,y) \in T} x + \max_{(x,y) \in T} y - \min_{(x,y) \in T} y.$$

A *Steiner tree* for  $T$  is a tree  $Y$  with  $T \subseteq V(Y) \subsetneq \mathbb{R}^2$ . We denote by  $\text{Steiner}(T)$  the length of a shortest rectilinear (i.e. edge lengths acc. to  $\ell_1$ ) Steiner tree for  $T$ . Moreover let  $\text{MST}(T)$  be the length of a minimum spanning tree in the complete graph on  $T$  with edge costs  $\ell_1$ .

Prove that:

- (a)  $\text{BB}(T) \leq \text{Steiner}(T) \leq \text{MST}(T)$ ;
- (b)  $\text{Steiner}(T) \leq \frac{3}{2} \text{BB}(T)$  for  $|T| \leq 5$ ;
- (c) There is no  $\alpha \in \mathbb{R}$  s.t.  $\text{Steiner}(T) \leq \alpha \text{BB}(T)$  for all finite  $\emptyset \neq T \subset \mathbb{R}^2$ .

(2 + 3 + 2 points)

**Exercise 4.2.** (a) Show that the Steiner ratio for the Euclidean plane is at least  $2/\sqrt{3}$ .

- (b) Show that the Steiner ratio for a metric space is at most  $2 - 2/n$  where  $n$  is the number of terminals.

(2 + 3 points)

**Exercise 4.3.** Consider the following algorithm to compute a rectilinear Steiner tree  $T$  for a set  $P$  of points in the plane  $\mathbb{R}^2$ .

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1: Choose  $p \in P$  arbitrarily;
2:  $T := (\{p\}, \emptyset)$ ,  $S := P \setminus \{p\}$ 
3: while  $S \neq \emptyset$  do
4:   Choose  $s \in S$  with minimum  $dist(s, T)$ 
5:   Let  $\{u, w\} \in E(T)$  be an edge which minimizes  $dist(s, SP(u, w))$ 
6:    $v := \arg \min\{dist(s, v) \mid v \in SP(u, w)\}$ 
7:    $T := (V(T) \cup \{v\} \cup \{s\}, E(T) \setminus (u, w) \cup \{u, v\} \cup \{v, w\} \cup \{v, s\})$ 
8:    $S := S \setminus \{s\}$ 
9: end while
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In this notation  $SP(u, w) \subset \mathbb{R}^2$  is the area covered by shortest paths between  $u$  and  $w$ , and  $dist(s, T)$  is the minimum distance between  $s$  and the shortest path area  $SP(u, w)$  of an edge  $\{u, w\} \in E(T)$ .

Show that the algorithm is a  $\frac{3}{2}$ -approximation algorithm for the MINIMUM STEINER TREE PROBLEM.

*Hint:* First show that the length of  $T$  is at most the length of a minimum spanning tree on  $P$ .

(8 points)

**Deadline:** **Thursday**, May 2<sup>nd</sup>, before the lecture. The websites for lecture and exercises can be found at:

<http://www.or.uni-bonn.de/lectures/ss19/chipss19.html>

In case of any questions feel free to contact me at [klotz@or.uni-bonn.de](mailto:klotz@or.uni-bonn.de).