

## Exercise Set 11

### Exercise 11.1:

Given a complete graph  $G$ ,  $s, t \in V(G)$  and weights  $c : E(G) \rightarrow \mathbb{R}_+$  satisfying the triangle inequality, consider the problem to find an  $s$ - $t$  path containing all vertices of  $G$  of minimum total weight.

Generalize CHRISTOFIDES' ALGORITHM to obtain a  $\frac{5}{3}$ -approximation.

(4 points)

### Exercise 11.2:

Let  $G$  be a complete undirected graph in which all edge lengths are either 1 or 2. Give a  $\frac{4}{3}$ -approximation algorithm for the TRAVELING SALESMAN PROBLEM in this special case.

*Hint: You may use that a minimum 2-matching in  $G$ , i.e. a subgraph of  $G$  in which every vertex has degree 2, can be computed in polynomial time.*

(4 points)

### Exercise 11.3:

Describe a polynomial-time algorithm which optimally solves any instance of the TRAVELING SALESMAN PROBLEM that is the metric closure of a weighted tree.

(4 points)

### Exercise 11.4:

The METRIC BIPARTITE TRAVELING SALESMAN PROBLEM is the problem of finding a Hamiltonian circuit of minimum cost in a bipartite graph  $H$  with a nonnegative cost function  $d$  satisfying

$$d(\{a, b\}) + d(\{a', b\}) + d(\{a', b'\}) \geq d(\{a, b'\}) \text{ for } \{a, b\}, \{a', b\}, \{a, b'\}, \{a', b'\} \in E(H).$$

Prove that for any  $k$ , if there is a  $k$ -factor approximation algorithm for the METRIC BIPARTITE TRAVELING SALESMAN PROBLEM, there is also a  $k$ -factor approximation algorithm for the METRIC TRAVELING SALESMAN PROBLEM.

*Hint: Given an instance  $(G, c)$  of the METRIC TSP, construct an instance  $(H, d)$  of the METRIC BIPARTITE TSP where  $V(H) := V(G) \times \{0, 1\}$  and  $d(\{(v, 0), (w, 1)\}) \in \{c(\{v, w\}), 0\}$ .*

(4 points)

**Deadline:** Thursday, July 13th, before the lecture.