

Exercise Set 8

Exercise 8.1:

Consider the PLACEMENT LEGALIZATION PROBLEM with $y_{\max} - y_{\min} = 1$. We are given an infeasible placement $\tilde{x} : \mathcal{C} \rightarrow \mathbb{R}$. Show that there are instances for which there is no optimum solution which is consistent with \tilde{x} , i.e. such that $x(C) < x(C') \Rightarrow \tilde{x}(C) \leq \tilde{x}(C')$. (4 points)

Exercise 8.2:

Consider the following variant of the SINGLE ROW PLACEMENT WITH FIXED ORDERING problem, in which we minimize the bounding box net length:

Input: A set $\mathcal{C} = \{C_1, \dots, C_n\}$ of circuits, widths $w(C_i) \in \mathbb{R}_+$, an interval $[0, w(\square)]$, s.t. $\sum_{i=1}^n w(C_i) \leq w(\square)$. A netlist $(\mathcal{C}, P, \gamma, \mathcal{N})$ where the offset of a pin $p \in P$ satisfies $x(p) \in [0, w(\gamma(p))]$. Weights $\alpha : \mathcal{N} \rightarrow \mathbb{R}_+$.

Task: Find a feasible placement given by a function $x : \mathcal{C} \rightarrow \mathbb{R}$ s.t. $0 \leq x(C_1)$, $x(C_i) + w(C_i) \leq x(C_{i+1})$ for $i = 1, \dots, n-1$ and $x(C_n) + w(C_n) \leq w(\square)$, that minimizes

$$\sum_{N \in \mathcal{N}} \alpha(N) \cdot \text{BB}(N).$$

Here, $\text{BB}(N)$ denotes the bounding box net length.

Show that there exist $f_i : [0, w(\square)] \rightarrow \mathbb{R}$, $i = 1, \dots, n$, piecewise linear, continuous and convex, such that we can solve this problem by means of the SINGLE ROW ALGORITHM. (4 points)

Exercise 8.3:

Give a detailed proof of Corollary 3.27, i.e. show that the variant of the PLACEMENT LEGALIZATION PROBLEM in which the goal is to minimize $\sum_{C \in \mathcal{C}} |x(C) - \tilde{x}(C)|$ is strongly NP-hard even if $\mathcal{S} = \emptyset$ and $y_{\max} - y_{\min} = 1$.

(3 points)

Exercise 8.4:

Programming exercise

Implement the dynamic programming buffering algorithm proposed by van Ginneken, with runtime $O(|L|^2|V(A)|^2)$. Your program will take one argument, the path to a file which contains the input, and print the result to the standard output. The task is to find a buffering that maximizes the worst slack. *Buffers can be placed only at vertices v with $|\delta^+(v)| = |\delta^-(v)| = 1$.*

Input: the first line of the input contains the number of vertices $n = |V(A)|$. The set of vertices is $\{0, \dots, n-1\}$. Each of the following $n-1$ lines contains four natural numbers $v w C_e R_e$, encoding an edge $e = (v, w)$ with capacitance C_e and resistance R_e . You may assume that this graph is an arborescence rooted at 0, in which every leaf is a sink.

Turn page! \rightarrow

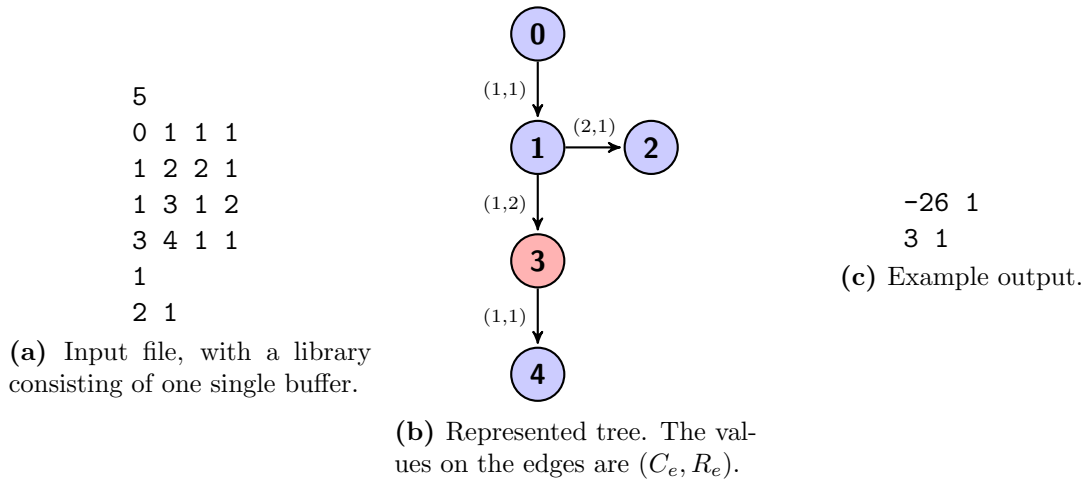


Figure 1: Example input and output, where a buffer of type 1 is placed at vertex 1. Buffers can be placed only at red vertices.

The next line contains the number of buffer types $|L|$ in the library. Each of the following lines contains two natural numbers $C_l R_l$, where C_l is the input capacitance and R_l the resistance of a buffer $l \in L$. The delay of a circuit (buffer or circuit at the root) with resistance R is $R \cdot dc$, where dc is the downstream capacitance.

For every sink v , assume $rat(v) = 0$, and input capacitance $C_v = 1$. Assume that the root has resistance $R_0 = 1$.

Output: the output must start with a line containing the worst slack and the number of buffers b . This should then be followed by b lines, consisting of two integers $v l$, which means that a buffer of type l is used at vertex v .

The program must be written in C or C++ (you are allowed to use up to C++11) and must compile and run on Linux. To achieve the maximum score, your program must not leak any memory and must be well documented. It must compile with either Clang $\geq 3.4.2$ or Gcc $\geq 4.8.3$, with `-Wall -Wpedantic -Werror`, and it cannot link to any other library. You can use any tool available in the standard library.

The deadline for the programming exercise is July 3rd, 12:00. Deliver your source code by email.

(20 points)

Deadline: June 23rd, before the lecture. The websites for lecture and exercises can be found at

<http://www.or.uni-bonn.de/lectures/ss16/ss16.html>

In case of any questions feel free to contact me at saccardi@or.uni-bonn.de.