

Exercise Sheet 9

Exercise 9.1:

Consider the EUCLIDEAN TRAVELING SALESMAN PROBLEM.

- (i) Prove that in an optimal solution of the problem no two edges cross each other.
- (ii) Find a class of instances for which the quality of the DOUBLE-TREE ALGORITHM is arbitrarily close to the guaranteed factor of 2.

(2 + 2 points)

Exercise 9.2:

Consider the following two formulations of the ASYMMETRIC TSP:

- (a) Given a strongly connected digraph G on n vertices with edge weights $c : E(G) \rightarrow \mathbb{R}_+$, find a minimum-cost connected multi-subgraph H such that $0 < |\delta_H^+(v)| = |\delta_H^-(v)|$ for all $v \in V(G)$, i. e. H is an Eulerian spanning multi-subgraph of G . A *multi-subgraph* arises from a subgraph by possibly duplicating edges.
- (b) Given a complete digraph G on n vertices with edge weights $c : E(G) \rightarrow \mathbb{R}_+$ satisfying the triangle inequality, i. e. $c((x, y)) + c((y, z)) \geq c((x, z))$ for all distinct $x, y, z \in V(G)$, find a minimum cost directed Hamiltonian circuit.

Prove the following:

- (i) Formulations (a) and (b) are equivalent, i. e. there are L-reductions in both directions with $\alpha = \beta = 1$.
- (ii) For every $n \geq 3$, there is an instance of (a) in which every optimum solution uses one of the edges at least $n - 2$ times.

(2 + 2 points)

Exercise 9.3:

Show that the EUCLIDEAN STEINER TREE PROBLEM has an approximation scheme.

Hint: Modify Arora's algorithm.

(4 points)

Exercise 9.4:

Given a complete graph G , $s, t \in V(G)$ and weights $c : E(G) \rightarrow \mathbb{R}_+$ satisfying the triangle inequality, consider the problem to find an s - t -path containing all vertices of G of minimum total weight.

- (i) Propose a 2-approximation algorithm for this problem.
- (ii) Generalize CHRISTOFIDES' ALGORITHM to obtain a $\frac{5}{3}$ -approximation algorithm.

(1 + 3 points)

Please hand in your solutions before the lecture on Tuesday, **June 17th, at 2:15 PM.**