

## Exercise Sheet 6

### Exercise 6.1:

Give a polynomial-time algorithm  $A$  for BIN PACKING such that there is a  $\delta > 0$  with  $A(I) \leq \text{OPT}(I) + \mathcal{O}(\text{OPT}(I)^{1-\delta})$  for all instances  $I$ .

(3 points)

### Exercise 6.2:

Let  $PCP'(\log n, 1)$  be defined as  $PCP(\log n, 1)$ , except that for each decision problem in  $PCP'(\log n, 1)$  there is a polynomial  $p$  such that the error probability (*soundness*) is at most  $1 - \frac{1}{p(\text{size}(x))}$  for input  $x$  instead of  $\frac{1}{2}$ . Prove that  $NP \subseteq PCP'(\log n, 1)$ . *Note:*

You may not use the  $PCP$ -Theorem.

(3 points)

### Exercise 6.3:

Prove that if  $\text{SATISFIABILITY} \in PCP(o(\log n), 1)$  (i. e.  $r \in \{0, 1\}^{o(\log(p(\text{size}(x))))}$ ) then  $P = NP$ .

(5 points)

### Exercise 6.4:

Consider the following problem: Given  $n$  variables  $x_1, \dots, x_n$  and a set of  $m$  boolean functions  $\Phi = \{\Phi_1, \dots, \Phi_m\}$  using  $k$  variables each, find a truth assignment such that the number of satisfied functions is maximized. Show that there is a constant  $k$  for which no polynomial-time 2-approximation algorithm for this problem exists (unless  $P = NP$ ).

(5 points)

Please hand in your solutions before the lecture on Tuesday, **May 20th, at 2:15 PM**.