

## Exercise Sheet 8

**Exercise 8.1:**

Describe a polytope  $P$  with a polynomial number of variables and constraints (a so-called *extended formulation*) such that the projection of  $P$  to some of the variables is the subtour polytope.

(4 points)

**Exercise 8.2:**

Let  $G = (V, E)$  be a 2-vertex-connected graph and  $(V, S)$  a DFS-tree in  $G$  rooted at  $r \in V$ . For each edge  $e = \{v, w\} \in E \setminus S$ , let w.l.o.g.  $v$  be on the  $r$ - $w$ -path in  $(V, S)$ , and let  $v'$  be the successor of  $v$  on this path. Add  $e$  to  $R$ ; moreover if  $|\delta(v)| \geq 3$  and  $e' = \{v, v'\}$  has not yet been added to  $R$ , then add also  $e'$  to  $R$  and  $\{e, e'\}$  to  $\mathcal{P}$  (cf. Figure 1). Prove that the resulting  $(R, \mathcal{P})$  is a *removable pairing* in  $G$ .

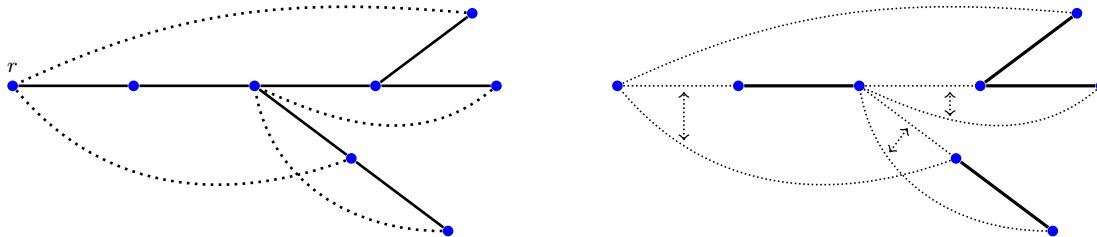


Figure 1: A 2-connected graph with a DFS tree (left: solid edges) and a removable pairing (right: dotted edges are in  $R$ ; arrows indicate pairs).

(4 points)

**Exercise 8.3:**

Consider the TRAVELING SALESMAN PROBLEM on subcubic graphs (i.e. graphs with maximum degree 3). Propose a  $\frac{4}{3}$ -approximation for this problem.

*Hint:* Use Exercise 8.2 and the Lemma of Mömke and Svensson.

(4 points)

**Exercise 8.4:**

Let  $G$  be a complete undirected graph in which all edge lengths are either 1 or 2. Give a  $\frac{4}{3}$ -approximation algorithm for the TRAVELING SALESMAN PROBLEM in this special case.

*Hint:* Find a minimum 2-matching in  $G$ , i.e. a subset  $M$  of edges such that every vertex has exactly 2 edges of  $M$  incident at it, and consider the cycles in  $M$ .

(4 points)

Please return your solutions before the lecture on Tuesday, **June 11th, 2:15 PM**.