

Exercise Sheet 3

Exercise 3.1:

An instance of MAXIMUM SATISFIABILITY is called k -satisfiable if any k of its clauses can be satisfied simultaneously. Let r_k be the infimum of the fraction of clauses that one can satisfy in any k -satisfiable instance. Show:

1. $r_1 = \frac{1}{2}$.
2. $r_2 = (\sqrt{5} - 1)/2$.

Hint: Set literals appearing in one-element clauses to *true* with probability a (for some $1/2 < a < 1$) and assign the remaining variables to *true* with probability $1/2$. Derandomize this and choose a appropriately.

(4 points)

Exercise 3.2:

The KNAPSACK PROBLEM for an instance \mathcal{I} can be formulated as integer program:

$$\max \left\{ \sum_{i=1}^n c_i x_i \mid \sum_{i=1}^n w_i x_i \leq W, x_i \in \{0, 1\} \forall 1 \leq i \leq n \right\} \quad (1)$$

Denote the optimum of (1) by $OPT(\mathcal{I})$ and let $LR(\mathcal{I})$ be the optimum of the linear relaxation where $x_i \in \{0, 1\}$ is replaced by $0 \leq x_i \leq 1$.

Show that the *integrality gap* $\sup \left\{ \frac{LR(\mathcal{I})}{OPT(\mathcal{I})} \mid OPT(\mathcal{I}) \neq 0 \right\}$ of the KNAPSACK PROBLEM is unbounded. Show that this doesn't remain valid if the definition of the KNAPSACK PROBLEM additionally requires $w_i \leq W$ for all w_i .

(4 points)

Exercise 3.3:

Consider the FRACTIONAL MULTI KNAPSACK PROBLEM: Given natural numbers $n, m \in \mathbb{N}$ and w_i, c_{ij} as well as W_j for $1 \leq i \leq n$ and $1 \leq j \leq m$, find x_{ij} satisfying $\sum_{j=1}^m x_{ij} = 1$ for all $1 \leq i \leq n$ and $\sum_{i=1}^n x_{ij} w_i \leq W_j$ for all $1 \leq j \leq m$ such that $\sum_{i=1}^n \sum_{j=1}^m x_{ij} c_{ij}$ is minimum.

State a polynomial-time combinatorial algorithm for this problem or prove NP -hardness.

(4 points)

Exercise 3.4:

Describe exact algorithms with running time $O(2^{n/2})$ for the following problems:

1. SUBSET SUM: Given $K, n, x_1, \dots, x_n \in \mathbb{N}$, find $S \subseteq \{1, \dots, n\}$ with $\sum_{i \in S} x_i = K$.
2. KNAPSACK PROBLEM (where n denotes the number of items).

(4 points)

Please return your solutions before the lecture at Tuesday, **April 30th, 2:15 PM**.