Winter semester 2013/14 Prof. Dr. S. Held

## Linear and Integer Optimization

## Exercise Sheet 9

**Exercise 9.1:** Show that for each  $K \in \mathbb{N}$  there is a bounded integer program with two variables such that the branch-&-bound algorithm visits more than K vertices in the branch-&-bound tree. (5 Points).

**Exercise 9.2:** Consider the following ILP (integer linear program)

$$\begin{array}{rl} \max -\sqrt{2}x + y \\ -\sqrt{2}x + y &\leq 0 \\ x &\geq 1 \\ y &\geq 0 \\ x, y &\in \mathbb{Z} \end{array}$$

Show:

- The ILP has feasible solutions. (1 Point)
  The objective is bounded from above by 0. (1 Point)
- 3. There is no optimum solution. (2 Points)
- 4. Let S be the set of solutions of the ILP, then conv(S) is not a polyhedron.

(1 Point)

**Exercise 9.3:** (Mixed integer programming hull) Let  $P = \{x \in \mathbb{R}^{k+l} : Ax \leq b\}$  be a rational polyhedron. Show that  $\operatorname{conv}(P \cap (\mathbb{Z}^k \times \mathbb{R}^l))$  is a rational polyhedron. (5 Points)

## Exercise 9.4:

The MAXIMUM-STABLE-SET-PROBLEM is defined as follows. Given a graph G, we are looking for a vertex set  $S \subseteq V(G)$  of maximum cardinality |S| such that  $\{v, w\} \notin E(G)$  for all  $v, w \in S$ . A vertex set S with  $E(G[S]) = \emptyset$  is called *stable set*. This problem can be modeled by the following ILP:

$$\max \sum_{v \in V(G)} x_v \tag{1}$$

s.d. 
$$x_v + x_w \le 1$$
  $\forall \{v, w\} \in E(G)$  (2)

$$x_v \in \{0, 1\} \qquad \qquad \forall v \in V(G) \tag{3}$$

Show:

1. There is a class of graphs  $(G_n)_{n \in \mathbb{N}}$  with  $|V(G_n)| = n$  such that the integrality gap has size  $\Omega(|V(G_n)|)$ , where (3) is replaced by

$$0 \le x_v \le 1 \quad \forall v \in V(G)$$

in the linear relaxation.

2. Following inequalities are valid for the MAXIMUM-STABLE-SET-PROBLEM:

$$\sum_{v \in V(C)} x_v \leq 1 \quad \text{for a clique } C \subseteq G,$$
$$\sum_{v \in V(H)} x_v \leq \left\lceil \frac{|V(H)| - 1}{2} \right\rceil \quad \text{for a circuit } H \subseteq G,$$

where a clique C is a complete subgraph of G, i.e.  $G[C] \simeq K_{|C|}$ . Give examples for both types of inequalities, where the optima of the linear relaxations are reduced (strictly).

(5 Points)

Submission deadline: Tuesday, 17.12.2013, before the lecture.