

## Combinatorial Optimization

### Exercise Sheet 3

#### Exercise 3.1:

Prove:

1. A minimal factor-critical graph  $G$  has at most  $\frac{3}{2}(|V(G)| - 1)$  edges. This bound is tight. (Needed for Lemma 1.34) (2 Points)
2. Let  $G$  be a graph and  $M$  a matching in  $G$ . If  $X \subseteq V(G)$  is the set of  $M$ -exposed vertices, then a shortest  $M$ -alternating  $X$ - $X$ -walk of positive length can be found in  $O(|E(G)|)$ . (Lemma 1.39) (2 Points)

#### Exercise 3.2:

Prove that an undirected graph  $G$  is factor-critical if and only if  $G$  is connected and  $\nu(G) = \nu(G - v)$  for all  $v \in V(G)$ . (3 Points)

#### Exercise 3.3:

Let  $G$  be a graph and  $M$  a matching in  $G$  that is not maximum.

1. Show that there are  $\nu(G) - |M|$  vertex-disjoint  $M$ -augmenting paths in  $G$ .  
*Hint:* Recall the proof of Berge's Theorem.
2. Prove that there exists an  $M$ -augmenting path of length at most  $\frac{\nu(G) + |M|}{\nu(G) - |M|}$ .
3. Let  $P$  be a shortest  $M$ -augmenting path in  $G$  and  $P'$  an  $(M \triangle E(P))$ -augmenting path. Prove  $|E(P')| \geq |E(P)| + |E(P \cap P')|$ .

Consider the following algorithm: We start with the empty matching and in each iteration augment the matching along a shortest augmenting path. Let  $P_1, P_2, \dots$  be the sequence of augmenting paths chosen.

4. Show that if  $|E(P_i)| = |E(P_j)|$  for  $i \neq j$ , then  $P_i$  and  $P_j$  are vertex-disjoint.
5. Conclude that the sequence  $|E(P_1)|, |E(P_2)|, \dots$  contains at most  $2\sqrt{\nu(G)} + 2$  different numbers. (5 Points)

**Exercise 3.4:**

Let  $G = (V, E)$  a graph and  $X \subseteq V$ . Let  $\beta(G, X)$  be the maximum size of a set  $Y \subseteq X$  for which there is a matching in  $G$  that covers  $Y$ . Prove

$$\beta(G, X) = \min_{U \subseteq V} \{|X| + |U| - q_X(U)\}.$$

Here  $q_X(U)$  denotes the number of odd connected components of  $(V \setminus U, E)$  whose vertices are all in  $X$ .

*Hint:* Construct a new graph with  $2|V|$  vertices and apply Tutte's Theorem. (4 Points)

**Deadline:** Tuesday, October 30, 2012, before the lecture.